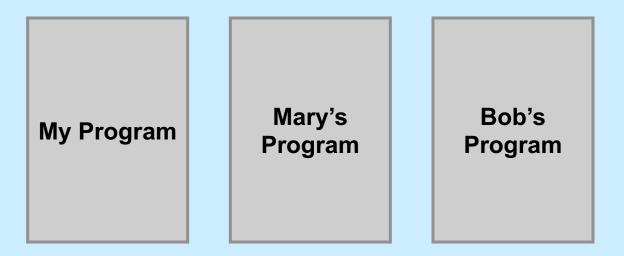
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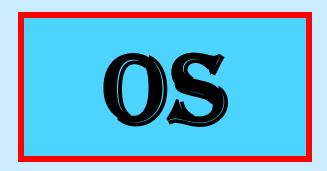
Architecture and the OS

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The Operating System





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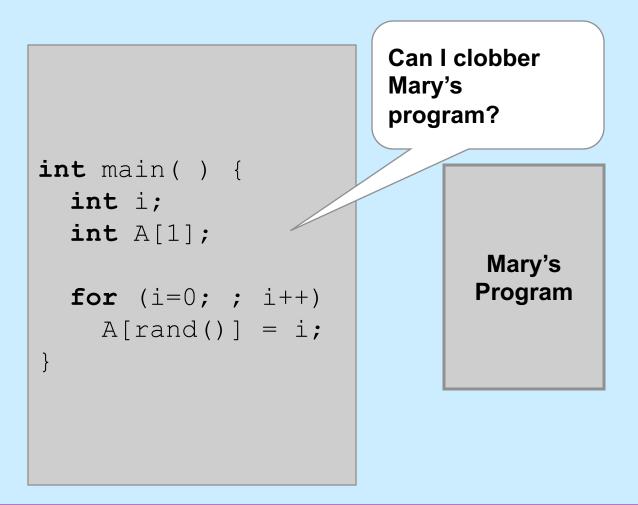
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Processes

Containers for programs

- virtual memory
 - » address space
- scheduling
 - » one or more threads of control
- file references
 - » open files
- and lots more!

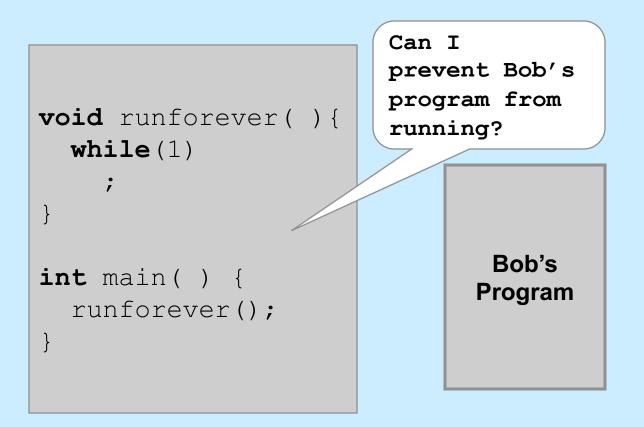
Idiot Proof ...



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Fair Share



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Architectural Support for the OS

- Not all instructions are created equal ...
 - non-privileged instructions
 - » can affect only current program
 - privileged instructions
 - » may affect entire system
- Processor mode
 - user mode
 - » can execute only non-privileged instructions
 - privileged mode
 - » can execute all instructions

Which Instructions Should Be Privileged?

- I/O instructions
- Those that affect how memory is mapped
- Halt instruction
- Some others ...

Who Is Privileged?

No one

user code always runs in user mode

- The operating-system kernel runs in privileged mode
 - nothing else does
 - not even super user on Unix or administrator on Windows

Entering Privileged Mode

- How is OS invoked?
 - very carefully ...
 - strictly in response to interrupts and exceptions
 - (booting is a special case)

Interrupts and Exceptions

- Things don't always go smoothly ...
 - I/O devices demand attention
 - timers expire
 - programs demand OS services
 - programs demand storage be made accessible
 - programs have problems
- Interrupts
 - demand for attention from external sources
- Exceptions
 - executing program requires attention

Exceptions

• Traps

- "intentional" exceptions
 - » execution of special instruction to invoke OS
- after servicing, execution resumes with next instruction
- Faults
 - a problem condition that is normally corrected
 - after servicing, instruction is re-tried
- Aborts
 - something went dreadfully wrong ...
 - not possible to re-try instruction, nor to go on to next instruction

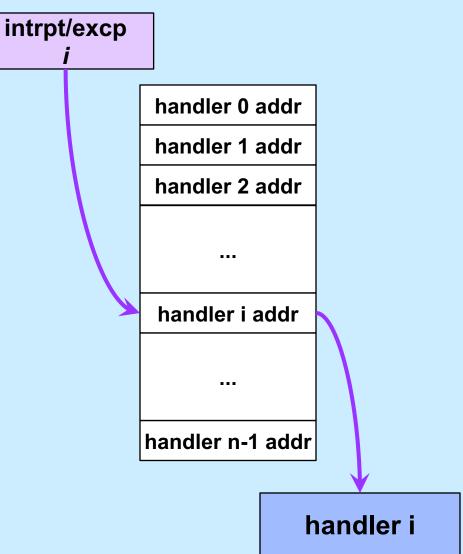
Actions for Interrupts and Exceptions

- When interrupt or exception occurs
 - processor saves state of current thread/process on stack
 - processor switches to privileged mode (if not already there)
 - invokes handler for interrupt/exception
 - if thread/process is to be resumed (typical action after interrupt)
 - » thread/process state is restored from stack
 - if thread/process is to re-execute current instruction
 - » thread/process state is restored, after backing up instruction pointer
 - if thread/process is to terminate
 - » it's terminated

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Interrupt and Exception Handlers

- Interrupt or exception invokes handler (in OS)
 - via interrupt and exception vector
 - » one entry for each possible interrupt/exception
 - contains
 - address of handler
 - code executed in privileged mode
 - » but code is part of the OS

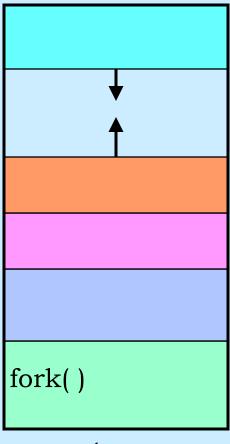


Creating Your Own Processes



```
#include <unistd.h>
int main() {
 pid t pid;
 if ((pid = fork()) == 0) {
      /* new process starts
         running here */
 /* old process continues
    here */
}
```

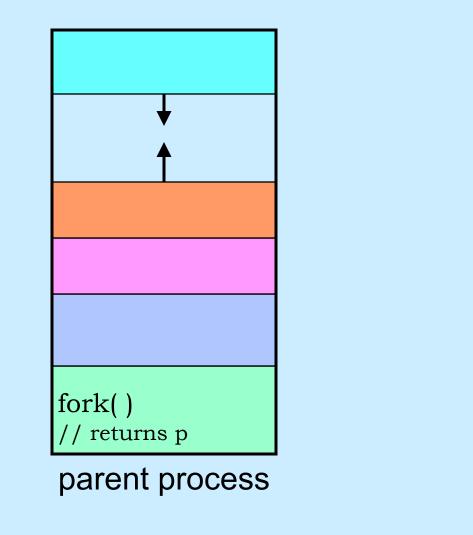
Creating a Process: Before

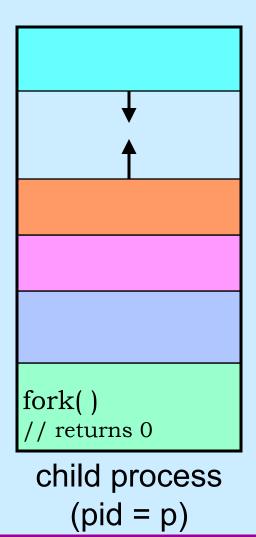


parent process

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Creating a Process: After





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Quiz 1

The following program

- a) runs forever
- b) terminates quickly

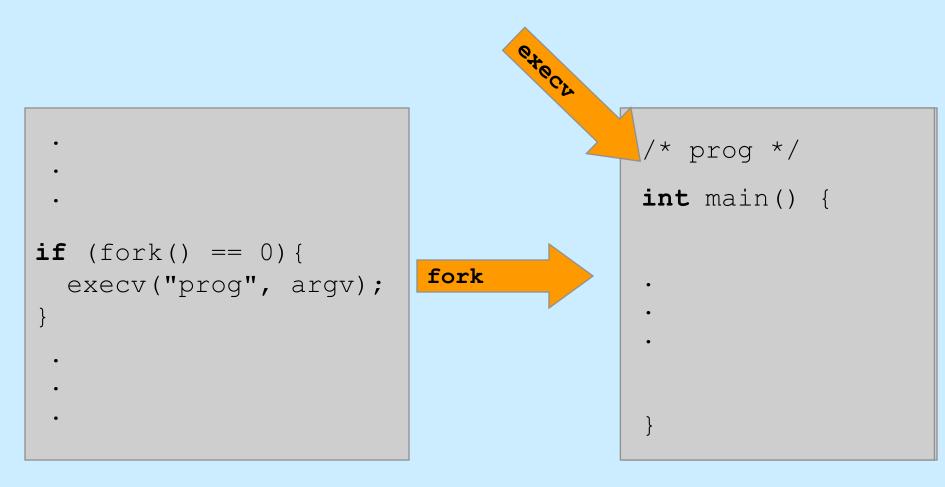
```
int flag;
int main() {
  while (flag == 0) {
    if (fork() == 0) {
      // in child process
      flag = 1;
      exit(0); // causes process to terminate
    }
```

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Process IDs

```
int main() {
                               parent prints:
 pid t pid;
                                 27355, 27342, 27342
 pid t ParentPid = getpid();
                               child prints:
 if ((pid = fork()) == 0) {
                                 0, 27342, 27355
      printf("%d, %d, %d\n",
            pid, ParentPid, getpid());
      return 0;
  }
 printf("%d, %d, %d\n",
            pid, ParentPid, getpid());
  return 0;
```

Putting Programs into Processes



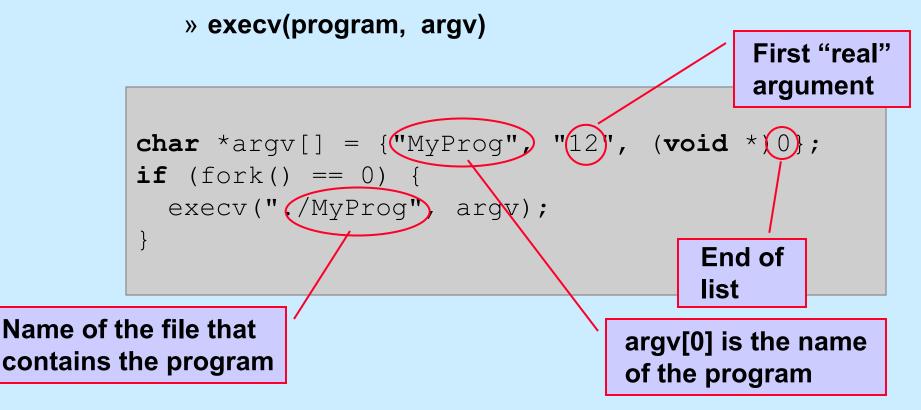
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Exec

Family of related system functions

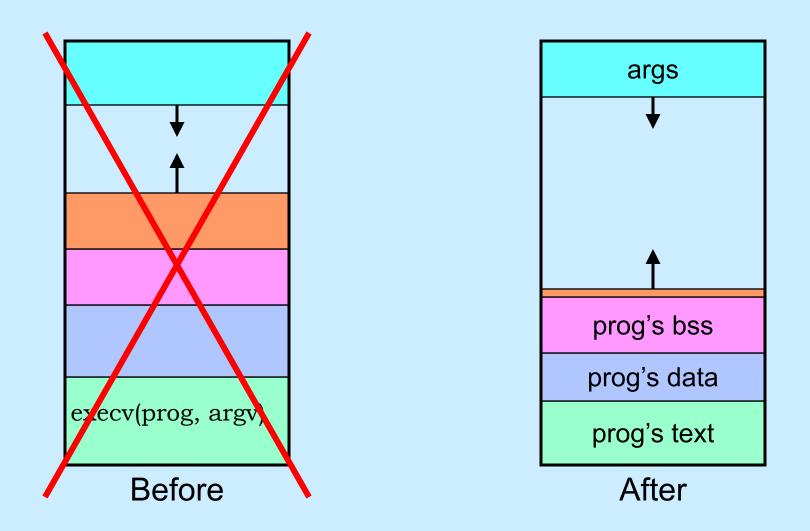
-we concentrate on one:



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Loading a New Image



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A Random Program ...

int main(int argc, char *argv[]) {

```
if (argc != 2) {
    fprintf(stderr, "Usage: random count\n");
    exit(1);
  }
  int stop = atoi(argv[1]);
```

```
for (int i = 0; i < stop; i++)</pre>
```

```
printf("%d\n", rand());
```

return 0;

Passing It Arguments

From the shell

\$ random 12

• From a C program

```
if (fork() == 0) {
    char *argv[] = {"random", "12", (void *)0};
    execv("./random", argv);
}
```

Quiz 2

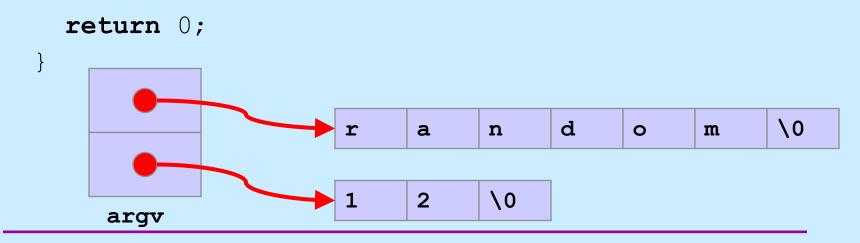
}

```
if (fork() == 0) {
    char *argv[] = {"random", "12", (void *)0};
    execv("./random", argv);
    printf("random done\n");
```

The *printf* statement will be executed a) always b) only if execv fails c) only if execv succeeds

Receiving Arguments

```
int main(int argc, char *argv[]) {
    if (argc != 2) {
        fprintf(stderr, "Usage: random count\n");
        exit(1);
    }
    int stop = atoi(argv[1]);
    for (int i = 0; i < stop; i++)
        printf("%d\n", rand());</pre>
```



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Not So Fast ...

How does the shell invoke your program?

if (fork() == 0) {
 char *argv = {"random", "12", (void *)0};
 execv("./random", argv);
}
/* what does the shell do here??? */

Wait

```
#include <unistd.h>
#include <sys/wait.h>
...
 pid t pid;
  int status;
  ...
  if ((pid = fork()) == 0) {
    char *argv[] = {"random", "12", (void *)0};
    execv("./random", argv);
  }
 waitpid(pid, &status, 0);
```

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Exit

```
#include <unistd.h>
#include <stdlib.h>
#include <sys/wait.h>
int main() {
 pid t pid;
  int status;
  if ((pid = fork()) == 0) {
    if (do work() == 1)
      exit(0); /* success! */
                                     exit code
    else
      exit(1); /* failure ... *
  }
 waitpid(pid, &status, 0);
  /* low-order byte of status contains exit code.
     WEXITSTATUS (status) extracts it */
```

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Shell: To Wait or Not To Wait ...

```
$ who
```

```
if ((pid = fork()) == 0) {
      char *argv[] = {"who", 0};
      execv("who", argv);
   }
   waitpid(pid, &status, 0);
   ...
Ŝ
 who &
   if ((pid = fork()) == 0) {
      char *argv[] = {"who", 0};
      execv("who", argv);
   }
```

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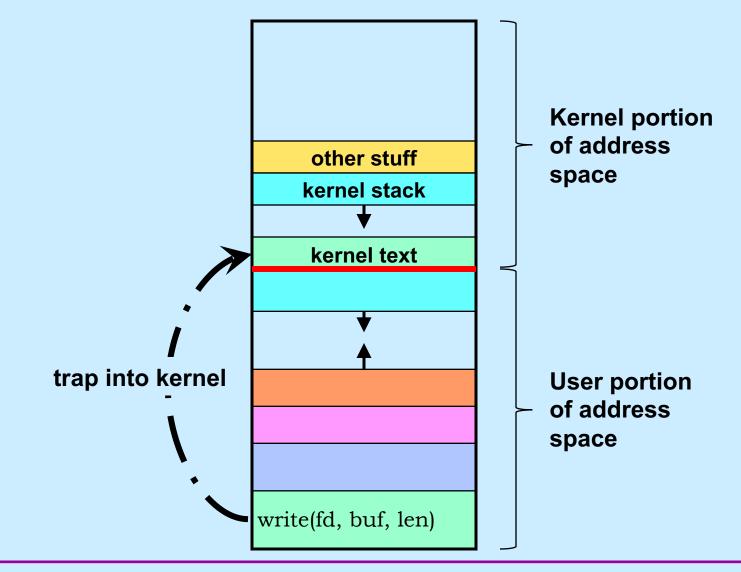
...

System Calls

- Sole direct interface between user and kernel
- Implemented as library functions that execute trap instructions to enter kernel
- Errors indicated by returns of –1; error code is in global variable errno

```
if (write(fd, buffer, bufsize) == -1) {
    // error!
    printf("error %d\n", errno);
    // see perror
}
```

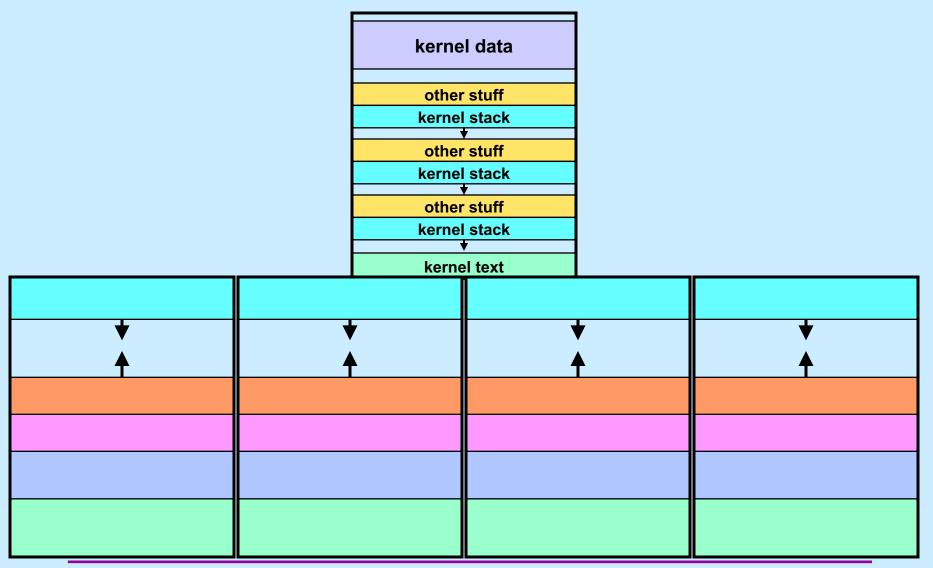
System Calls



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Multiple Processes



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Shells and Files

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Shells



- Command and scripting languages for Unix
- First shell: Thompson shell
 - sh, developed by Ken Thompson
 - released in 1971
- Bourne shell
 - also sh, developed by Steve Bourne
 - released in 1977
- C shell
 - csh, developed by Bill Joy
 - released in 1978
 - tcsh, improved version by Ken Greer

More Shells



- Bourne-Again Shell
 - bash, developed by Brian Fox
 - released in 1989
 - found to have a serious security-related bug in 2014
 - » shellshock
- Almquist Shell
 - ash, developed by Kenneth Almquist
 - released in 1989
 - similar to bash
 - dash (debian ash) used for scripts in Debian Linux
 - » faster than bash
 - » less susceptible to shellshock vulnerability

Roadmap

- We explore the file abstraction
 - what are files
 - how do you use them
 - how does the OS represent them
- We explore the shell
 - how does it launch programs
 - how does it connect programs with files
 - how does it control running programs

shell 2

shell 1

The File Abstraction

- A file is a simple array of bytes
- A file is made larger by writing beyond its current end
- Files are named by paths in a naming tree
- System calls on files are synchronous
- Files are permanent

Naming

- (almost) everything has a path name
 - files
 - directories
 - devices (known as special files)
 - » keyboards
 - » displays
 - » disks
 - » etc.

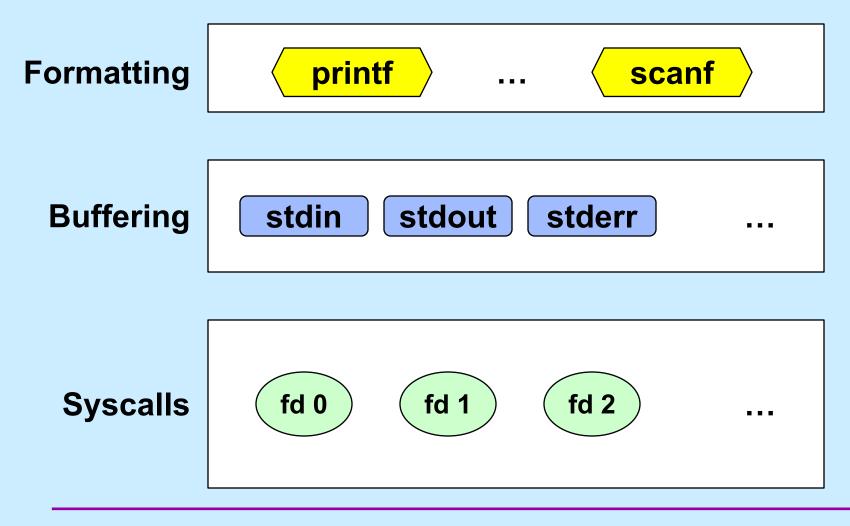
I/O System Calls

- int file_descriptor = open(pathname, mode [, permissions])
- int close (file descriptor)
- ssize_t count = read(file_descriptor, buffer_address, buffer_size)
- ssize_t count = write(file_descriptor, buffer_address, buffer_size)

Standard File Descriptors

```
int main() {
 char buf[BUFSIZE];
 int n;
 const char *note = "Write failed\n";
 while ((n = read(0, buf, sizeof(buf))) > 0)
  if (write(1, buf, n) != n) {
         write(2, note, strlen(note));
         exit(1);
 return(0);
```

Standard I/O Library



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Standard I/O

FILE *stdin;

FILE *stdout;

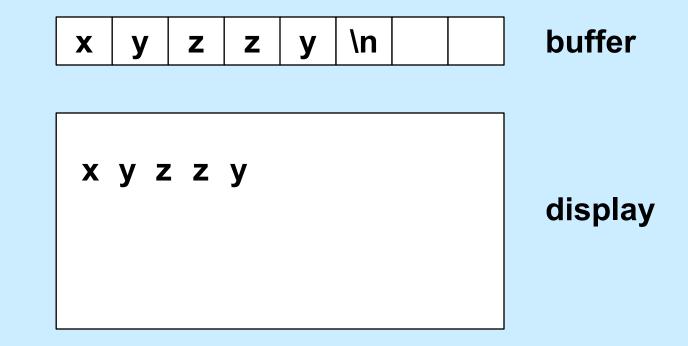
FILE *stderr;

// declared in stdio.h
// declared in stdio.h
// declared in stdio.h

Buffered Output

- printf("xy");
- printf("zz");

printf("y\n");



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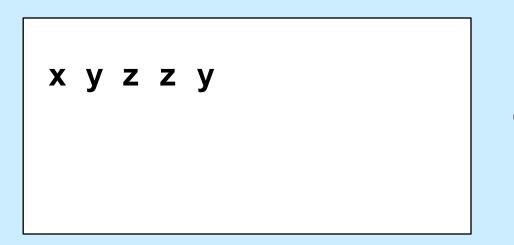
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Unbuffered Output

fprintf(stderr, "xy");

fprintf(stderr, "zz");

fprintf(stderr, "y\n");



display

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