CS 33

Multithreaded Programming VII

Implementing Mutexes

Strategy

- make the usual case (no waiting) very fast
- can afford to take more time for the other case (waiting for the mutex)

Futexes

- Safe, efficient kernel conditional queueing in Linux
- All operations performed atomically

- » otherwise return
- futex wake(futex t *futex)
 - » wake up one thread from futex's wait queue, if there are any waiting threads

Ancillary Functions

```
• int atomic inc(int *val)

    add 1 to *val, return its original value

• int atomic dec(int *val)

    subtract 1 from *val, return its original value

• int CAS(int *ptr, int old, int new) {
      int tmp = *ptr;
      if (*ptr == old)
          *ptr = new;
      return tmp;
```

Attempt 1

```
void lock(futex_t *futex) {
  int c;
  while ((c = atomic_inc(&futex->val)) != 0)
    futex_wait(futex, c+1);
}

void unlock(futex_t *futex) {
  futex->val = 0;
  futex_wake(futex);
}
```

Quiz 1

```
void lock(futex t *futex) {
  int c;
  while ((c = atomic inc(&futex->val)) != 0)
    futex wait(futex, c+1);
void unlock(futex t *futex)
  futex->val = 0;
  futex wake(futex);
```

Why doesn't Attempt 1 work?

- a) unlock fails to wake up a sleeping thread in certain circumstances
- b) the while loop in lock doesn't terminate in certain circumstances
- c) both of the above
- d) none of the above

Attempt 2

```
void lock(futex t *futex) {
  int c;
  if ((c = CAS(\&futex->val, 0, 1) != 0)
    do {
      if (c == 2 || (CAS(&futex->val, 1, 2) != 0))
        futex wait(futex, 2);
    while ((c = CAS(&futex->val, 0, 2)) != 0))
void unlock(futex t *futex) {
  if (atomic dec(&futex->val) != 1) {
    futex->val = 0;
    futex wake (futex);
```

Quiz 2

Does it work?

- a) always
- b) except for pathological cases
- rarely
- never

Memory Allocation

Multiple threads

One heap

Bottleneck?

Solution 0:

Use your malloc implementation but use mutexes to make it thread-safe

- 1) Use a single mutex to protect the heap
 - no concurrent access
- 2) Use a mutex per block
 - concurrent access to the heap

Quiz 3

Solution 0.2 is not used because

- a) Since the free list is circular, deadlock cannot be avoided
- b) Since each core accesses memory via its private L1 cache, memory blocks used by one thread cannot be safely shared with others
- c) There will be too many calls to lock and unlock mutexes, slowly things down a lot
- d) Since there must be a mutex per block, too much memory is wasted
- e) Something else

Not a Quiz

How can it be done better?

Solution 1

- Divvy up the heap among the threads
 - each thread has its own heap
 - no mutexes required
 - no bottleneck
- How much heap does each thread get?
- What if one thread frees memory malloc'd by another?

Solution 2

- Multiple "arenas"
 - each with its own mutex
 - thread allocates from the first one it can find whose mutex was unlocked
 - » if none, then creates new one
 - deallocations go back to original arena

Solution 3

- Global heap plus per-thread heaps
 - threads pull storage from global heap only when needed
 - freed storage goes to per-thread heap
 - » unless things are imbalanced
 - then thread moves storage back to global heap
 - mutex on only the global heap
- What if one thread frees memory malloc'd by another?

Malloc/Free Implementations

- ptmalloc
 - based on solution 2
 - in glibc (i.e., used by default)
- tcmalloc
 - based on solution 3
 - from Google
- Which is best?

Test Program

```
const unsigned int N=64, nthreads=32, iters=10000000;
int main() {
  void *tfunc(void *);
 pthread t thread[nthreads];
  for (int i=0; i<nthreads; i++) {
    pthread create(&thread[i], 0, tfunc, (void *)i);
    pthread detach(thread[i]);
 pthread exit(0);
void *tfunc(void *arg) {
  long i;
  for (i=0; i<iters; i++) {
    long *p = (long *) malloc(sizeof(long) * ((i%N) +1));
    free(p);
  return 0;
```

Not a Quiz

Which is fastest?

- a) glibc (i.e., standard Linux)
- b) Google

Compiling It ...

```
% gcc -o ptalloc alloc.c -lpthread
% gcc -o tcalloc alloc.c -lpthread -ltcmalloc
```

Running It (2014) ...

```
$ time ./ptalloc
real 0m5.142s
user 0m20.501s
sys 0m0.024s
$ time ./tcalloc
real 0m1.889s
user 0m7.492s
sys 0m0.008s
```

Running It (2024) ...

```
$ time ./ptalloc real 0m0.558s user 0m6.141s sys 0m0.020s $ time ./tcalloc real 0m0.400s user 0m4.458s sys 0m0.008s
```

What's Going On (2014)?

```
$ strace -c -f ./ptalloc
% time seconds usecs/call calls errors syscall
100.00 0.040002 13 3007 520 futex
$ strace -c -f ./tcalloc
% time seconds usecs/call calls errors syscall
 0.00 0.000000
                0 59 13 futex
```

What's Going On (2024)?

```
$ strace -c -f ./ptalloc
% time seconds usecs/call calls errors syscall
0.00 0.000000 0
                               futex
$ strace -c -f ./tcalloc
% time seconds usecs/call calls errors syscall
 0.38 0.000016 1 10
                                futex
```

```
#define N 64
#define npairs 16
#define allocsPerIter 1024
const long iters = 8*1024*1024/allocsPerIter;
#define BufSize 10240
typedef struct buffer {
  int *buf[BufSize];
  unsigned int nextin;
  unsigned int nextout;
  sem t empty;
  sem t occupied;
 pthread t pthread;
 pthread t cthread;
} buffer t;
```

```
int main() {
  long i;
 buffer t b[npairs];
  for (i=0; i<npairs; i++) {
   b[i].nextin = 0;
   b[i].nextout = 0;
    sem init(&b[i].empty, 0, BufSize/allocsPerIter);
    sem init(&b[i].occupied, 0, 0);
   pthread create(&b[i].pthread, 0, prod, &b[i]);
   pthread create(&b[i].cthread, 0, cons, &b[i]);
  for (i=0; i<npairs; i++) {
   pthread join(b[i].pthread, 0);
   pthread join(b[i].cthread, 0);
  return 0;
```

```
void *prod(void *arg) {
  long i, j;
  buffer t *b = (buffer t *) arg;
  for (i = 0; i<iters; i++) {
    sem wait(&b->empty);
    for (j = 0; j<allocsPerIter; j++) {</pre>
      b->buf[b->nextin] = malloc(sizeof(int)*((j%N)+1));
      if (++b->nextin >= BufSize)
       b->nextin = 0:
    sem post(&b->occupied);
  return 0;
```

```
void *cons(void *arg) {
  long i, j;
  buffer t *b = (buffer t *) arg;
  for (i = 0; i<iters; i++) {
    sem wait(&b->occupied);
    for (j = 0; j<allocsPerIter; j++) {</pre>
      free(b->buf[b->nextout]);
      if (++b->nextout >= BufSize)
       b->nextout = 0;
    sem post(&b->empty);
  return 0;
```

Running It (2014) ...

```
$ time ./ptalloc2
real 0m1.087s
user 0m3.744s
sys 0m0.204s
$ time ./tcalloc2
real 0m3.535s
user 0m11.361s
sys 0m2.112s
```

Running It (2024) ...

```
$ time ./ptalloc2
real 0m1.594s
user 0m8.778s
sys 0m2.551s
$ time ./tcalloc2
real 0m7.089s
user 0m59.871s
sys 0m11.220s
```

What's Going On (2014)?

```
$ strace -c -f ./ptalloc2
% time seconds usecs/call calls errors syscall
93.04 8.246196 117 70173 20775 futex
$ strace -c -f ./tcalloc2
% time seconds usecs/call calls errors syscall
99.92 47.796676 153 311012 7244 futex
```

What's Going On (2024)?